

Data Flow Testing

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Outline

From Paths to DU-Paths

DU Path Testing

Path Testing: Pros and Cons

Pros:

- ▶ Abstract **adequacy** criterion
- ▶ A measure of **testability**: program complexity

Cons:

- ▶ Too abstract: not clear **when** to use **which** criterion
- ▶ No use **data flow** and variable **dependencies**
- ▶ No use of **specification**

Flow Graphs

Defining nodes

Abstract models of program (control) structure.

Assume for the rest:

- ▶ a single **start** (entry) node: n_s ,
- ▶ a single **termination** (exit) node n_t ,
- ▶ one component, no dead code

Structured Loops

- ▶ a loop L is a strongly connected component
- ▶ entry node of a loop L is node $n \notin nodes(L)$, such that $n \rightarrow m$ for some $m \in nodes(L)$

Structured Loops

- ▶ a loop L is a strongly connected component
- ▶ entry node of a loop L is node $n \notin nodes(L)$, such that $n \rightarrow m$ for some $m \in nodes(L)$
- ▶ **exit** node of a loop L is node $n \notin nodes(L)$, such that $m \rightarrow n$ for some $m \in nodes(L)$
- ▶ **structured** loops: **unique entry** and **exit** nodes

Annotated Flow Graphs

Defining nodes

$DEF(n, v)$ holds (for a var. v and a node n), when n defines v .

Examples:

- ▶ $input(v)$, or
- ▶ $v := exp$

$$DEF(n) = \{v \mid DEF(n, v)\}$$

Annotated Flow Graphs

Using nodes

$USE(n, v)$ holds (for a var. v and a node n), when n uses the values of v . Examples:

- ▶ $output(v)$,
- ▶ $x := exp(v)$,
- ▶ *if* $cond(v)$ *then*, or
- ▶ *while* $cond(v)$ *do*, ...

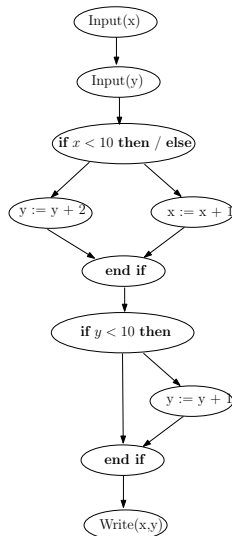
$$USE(n) = \{v \mid USE(n, v)\}$$

Also $REF(n, v)$ in the literature

Definitions and Uses: An Example

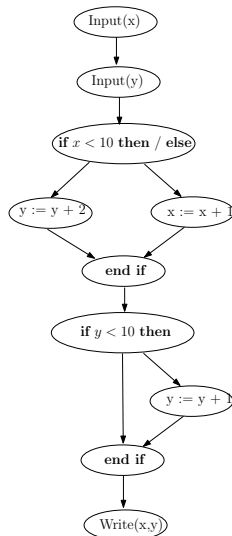
```

1: Input(x)
2: Input(y)
3: if  $x < 10$  then
4:    $y := y + 2$ 
5: else
6:    $x := x + 1$ 
7: end if
8: if  $y > 20$  then
9:    $y := y + 1$ ;
10: end if
11: Write(x,y)
12: end
  
```



Definitions and Uses: An Example

- 1: Input(x) {DEF(1) = {x}}
- 2: Input(y) {DEF(2) = {y}}
- 3: **if** $x < 10$ **then**
- 4: $y := y + 2$ {DEF(4) = USE(4) = {y}}
- 5: **else**
- 6: $x := x + 1$
- 7: **end if**
- 8: **if** $y > 20$ **then**
- 9: $y := y + 1$;
- 10: **end if**
- 11: Write(x,y) {USE(11) = {x,y}}
- 12: **end**



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From Paths to DU-Paths

DU Path Testing

DC Paths

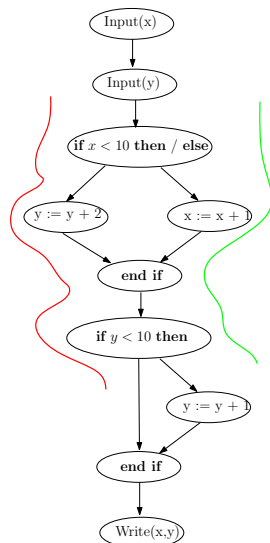
A definition-clear path (**DC-path**) wrt. v is a path m, P, m' such that for all $n \in \text{nodes}(P)$, $v \notin \text{DEF}(n)$.

$\text{DEF}(m, v)$ **reaches** $\text{USE}(n, v)$ when there is **DC-path** m, \dots, n .

DC Paths

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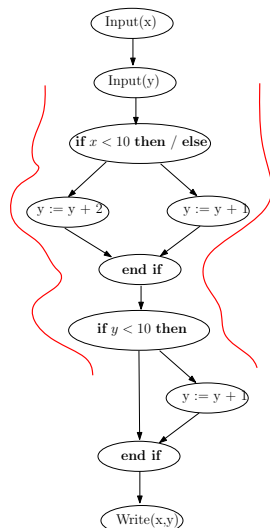
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DC Paths

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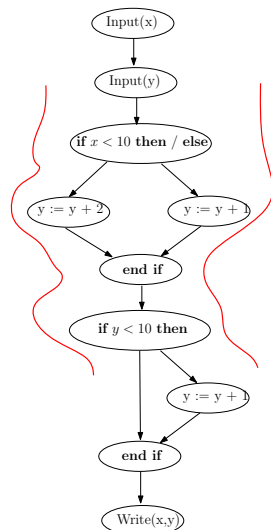
$\text{DEF}(m, v)$ **reaches** $\text{USE}(n, v)$ when there is **DC-path** m, \dots, n .



Anomalies

Unused definition: A $DEF(m, v)$ that does not reach any $USE(n, v)$.

Undefined usages: A DC-path wrt. v from n_s to n such that $USE(n, v)$.

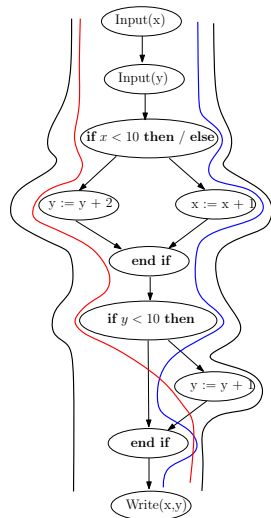


DU-Path Testing

A test-set S satisfies **DU-path** when for each variable v , for each nodes m, n, n' such that

1. $DEF(m, v)$ and $USE(n, v)$,
2. $DEF(m, v)$ reaches $USE(n, v)$, and
3. $n \rightarrow n'$

then **all** simple DC-paths (wrt. v) m, \dots, n, n' are covered by a test-case in S .



All-Uses Testing

A test-set S satisfies **All Uses** when
for each variable v , for each nodes m, n, n' such that

1. for each $DEF(m, v)$ and $USE(n, v)$,
2. $DEF(m, v)$ reaches $USE(n, v)$, and
3. $n \rightarrow n'$

then **at least one** simple DC-paths (wrt. v) m, \dots, n, n'
is covered by a test-case in S .

Def-Paths

For each variable v and node m such that $DEF(m, v)$, the Def-Path set of $DEF(m, v)$, denoted by $Def - Paths(m, v)$ is the set of all DC-paths m, \dots, n, n' such that:

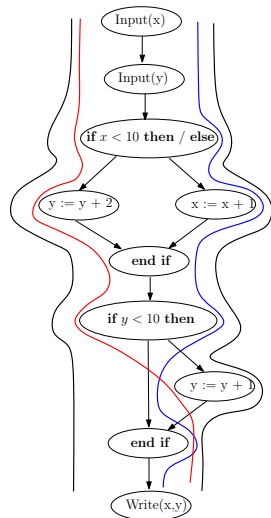
1. $USE(n, v)$,
2. $DEF(m, v)$ reaches $USE(n, v)$, and
3. $n \rightarrow n'$

All-Defs Path Testing

A test-set S satisfies **All-Defs** when for each variable v , for each nodes m, n, n' such that

1. for each $DEF(m, v)$, $USE(n, v)$,
2. for each path p in $Def - Paths(n, v)$,

there is a test case in S covering p .



Subsumption Relation

