

# Generative aspects in skeletal systems

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GRID.it  
project



Core GRID



# Summary

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- Skeletons
- Classical skeleton implementation (template)
- Alternative implementation (mdf + rewrules)
- Layered implementation
- Conclusions



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# Skeletons

- Useful, parametric, efficient parallelism exploitation pattern
  - **useful** : for a large class of applications
  - **parametric** : in the seq code, parallelism degree, types of tasks and results
  - **efficient** : known efficient implementations on a range of architectures



# Sample skel: farm

- **farm**: ('a → 'b) → stream 'a → stream 'b  
 $\text{farm } f \langle x_1, \dots, x_n \rangle = \langle f(x_1), \dots, f(x_n) \rangle$   
 $f(x_i)$  independent of  $f(x_j)$  any  $i, j$  (parallel)
- useful most of currently large scale parallel application fit the schema
- parametric in code, data types and parallelism degree
- efficient master slave, SMP multithread, ...



# Pros

- Separation of concerns
  - user: qualitative aspects
  - system: quantitative ones + mechanisms
- Portability (source code)
- Optimizations
  - exploit high level (meta) info
  - e.g. comm clustering, source rewriting, ...



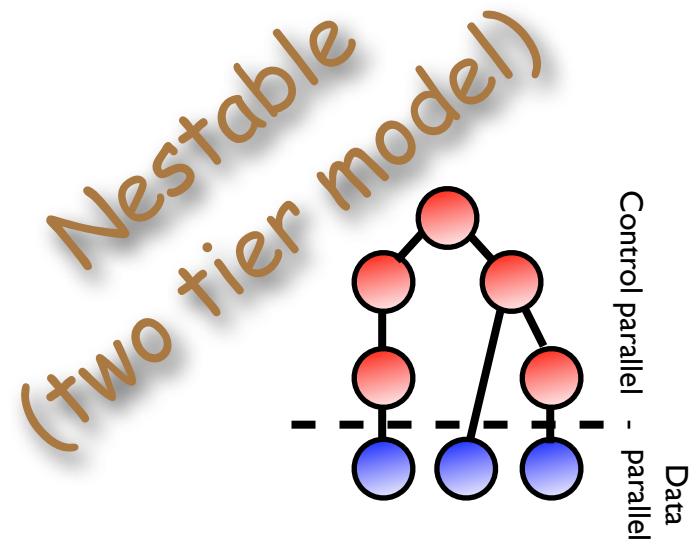
# Cons

- fixed skeleton set
  - impossibility to exploit slightly different parallel patterns
  - impossibility to introduce new parallel patterns
- poor / no interoperability with other parallel frameworks



# Typical skeletons

- Control parallel (aka stream parallel):
  - pipelines, task farms, loops, ...
  - parallelism : computation of different stream items
- Data parallel
  - map, reduce, prefix, ...
  - parallelism : computation of different (possibly overlapping) partitions of the same stream item





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# Generative aspects

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- Very diverse formalisms
  - skeletons : high level coordination patterns
  - object code : sequential language plus communication library calls
- Exploit existing knowledge and heuristics
  - to produce efficient, realistic obj code (template based implementations)



# Template based impl

```

seq f in(...) out(...)
$c{ ... }c$

seq g in(...) out(...)
$c{ ... }c$

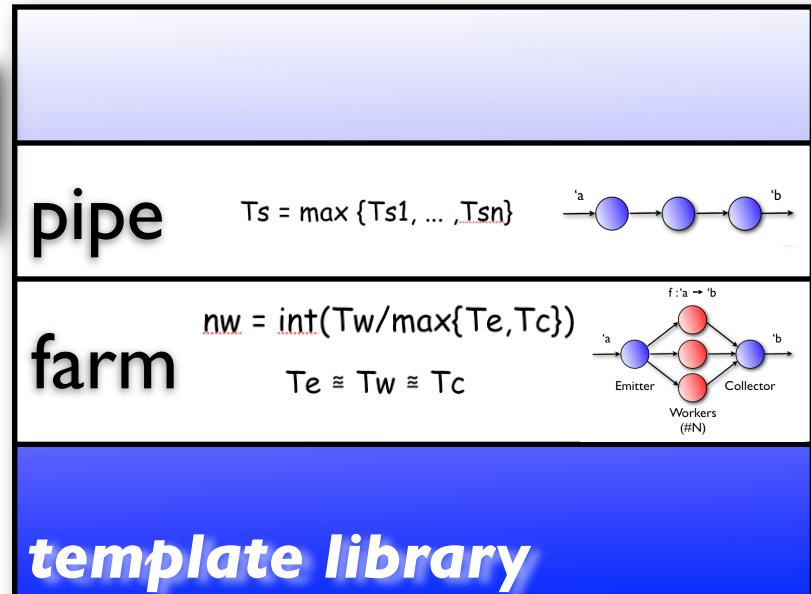
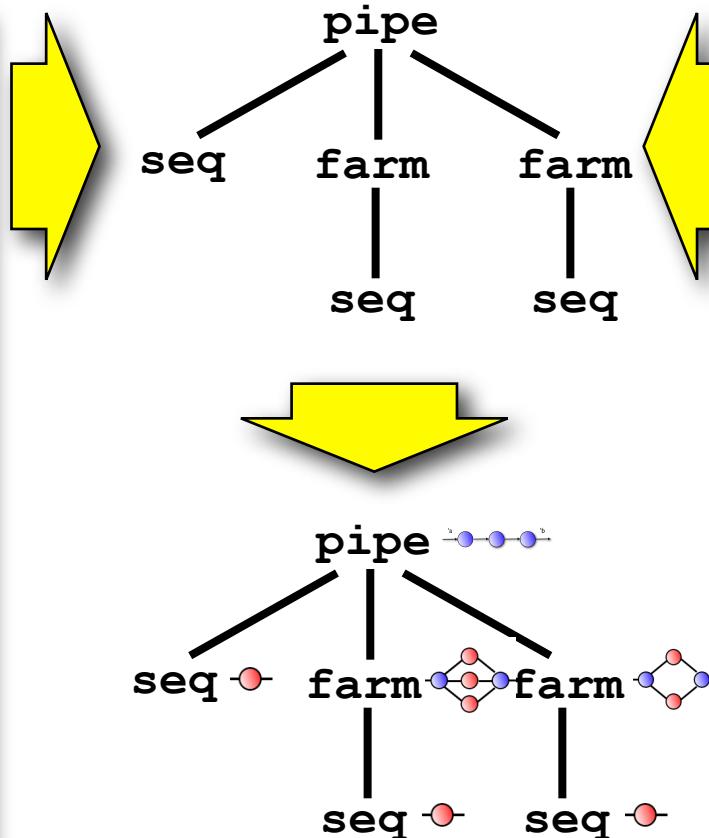
seq h in(...) out(...)
$c{ ... }c$

farm s2 in(...) out(...)
  g in(...) out(...)
end farm

farm s3 in(...) out(...)
  h in(...) out(...)
end farm

pipe main in(...) out(...)
  f in(...) out(...)
  s2 in(...) out(...)
  s3 in(...) out(...)
end pipe
  
```

**source code**

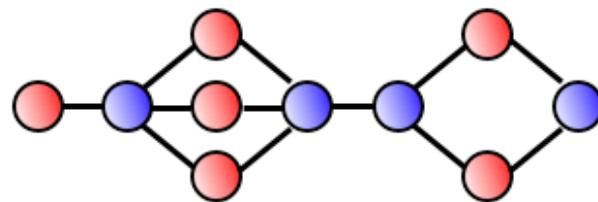


this is P3L ('92)

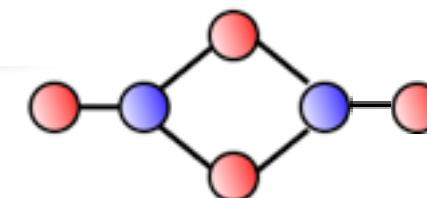
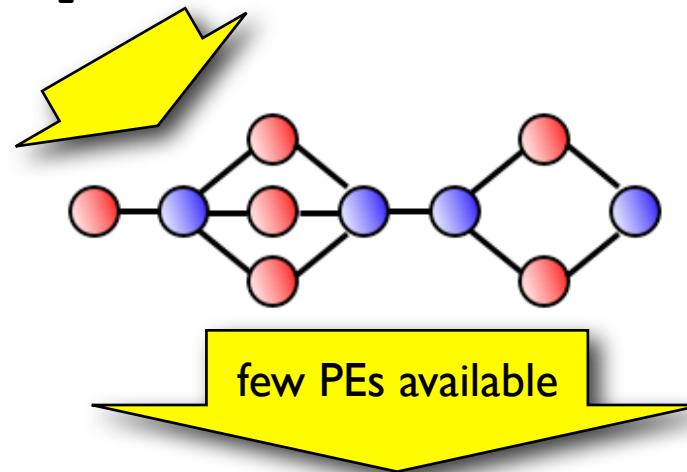
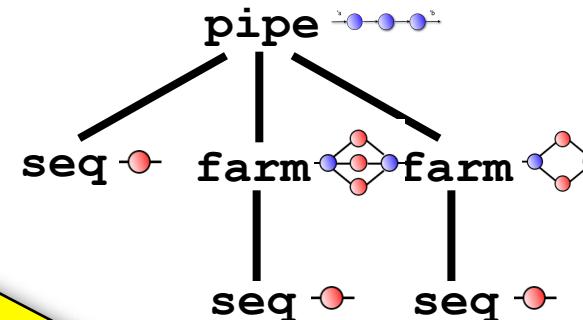
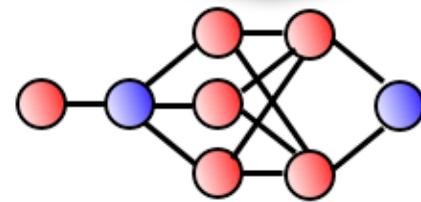
# Template based impl



Optimizations



$$(\alpha f) o (\alpha g) = \alpha(f o g)$$



Adaptivity



# Problems & issues

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- proper mechanisms to implement templates
  - decided at compile time (target nodes?)
- performance models
  - need rough estimate of run/comm times  
(profiling? user supplied?)
- optimizations
  - need knowledge relative to target config



# Template libs

- Template based, run time only implementations
- Library calls (user responsibility) to: define skel tree, invoke execution
- Library call implementation:
  - de facto instantiate templates
  - with additional parameters from users

this is Muesly ('02)



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# Generative aspects



- First skel generation (P3L, Muesli, eSkel)
  - code generated is virtually exposed to programmers (they do know its structure)
- Then: source2source transforms, optims, non template based implement. techniques
  - programmers unaware of what's going on
  - skeletons → meta info to be exploited

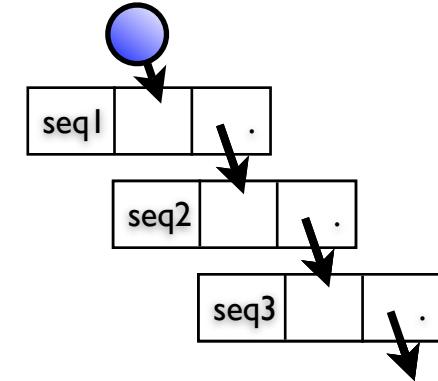
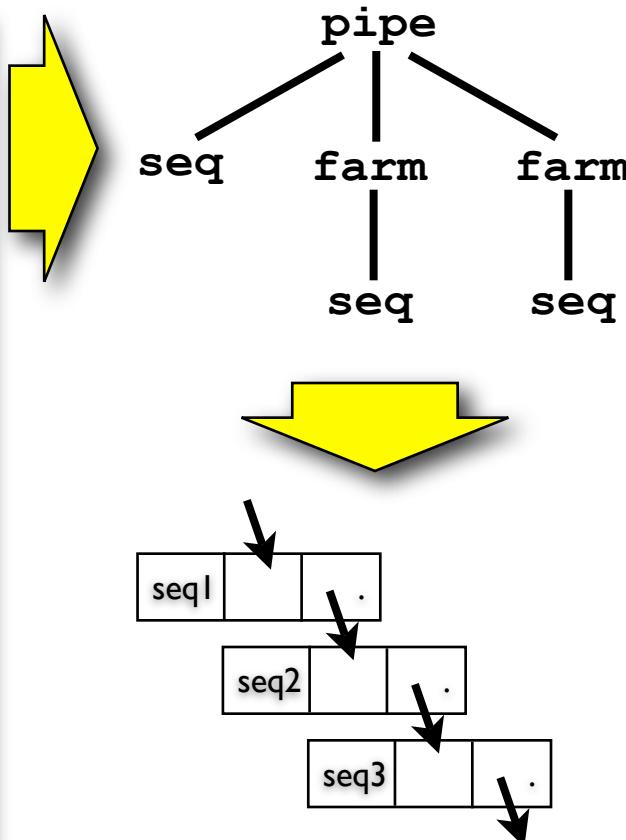
# Macro data flow impl



```
import muskel.*;
public class Main {
public static void
main(String[] args) {

Compute seq1 = new Seq1();
Compute seq2 = new Seq2();
Compute seq3 = new Seq3();
Farm s2 = new Farm(seq2);
Farm s3 = new Farm(seq3);
Pipeline farms =
new Pipeline(s2,s3);
Pipeline main =
new Pipeline(seq1,farms);
Manager mng =
new Manager(main);
m.inputStream("in.dat");
m.outputStream("o.dat");
m.compute();
return;
}
```

source code



this is Lithium/muskel ('00)



# Macro data flow impl

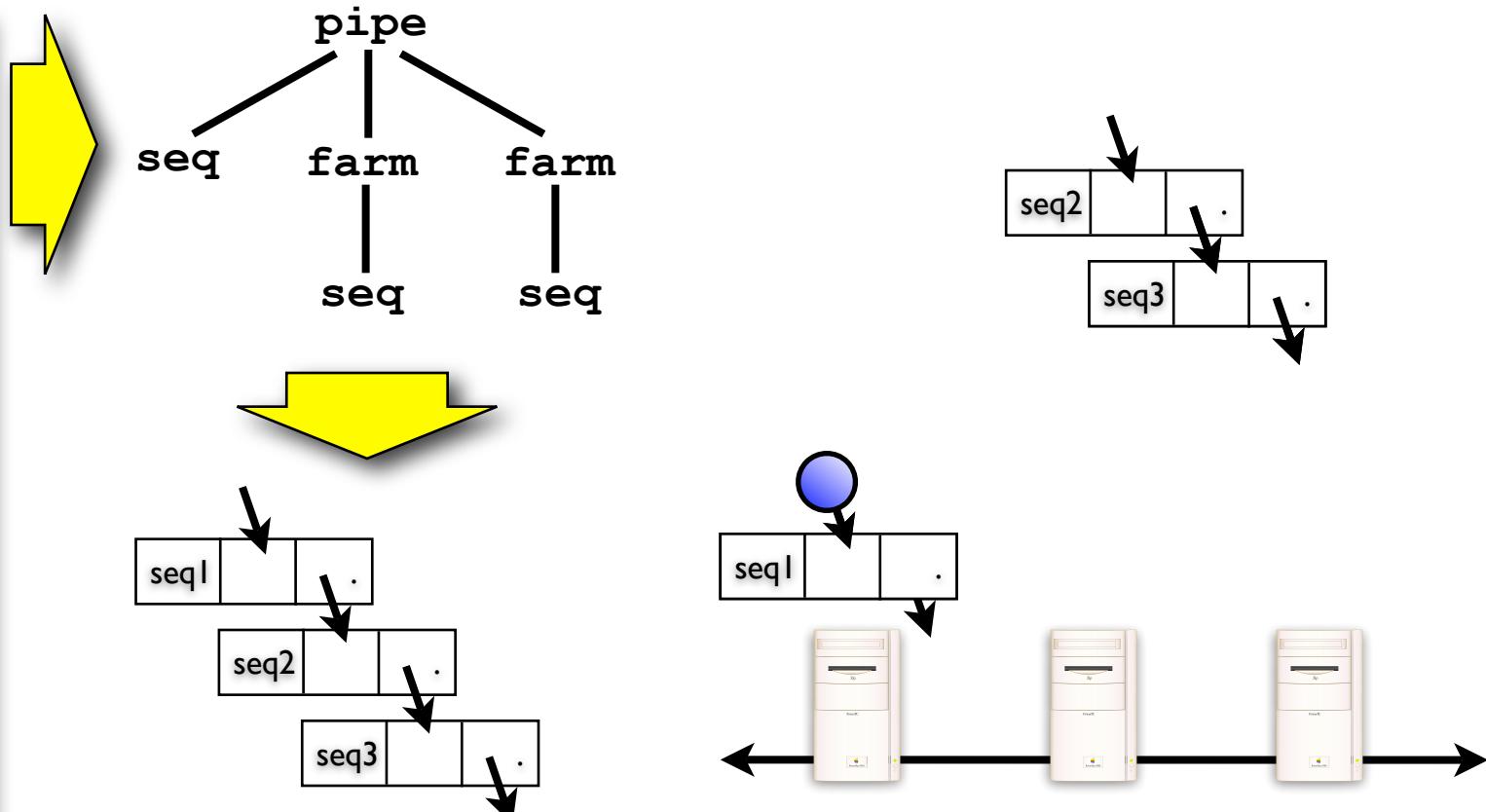
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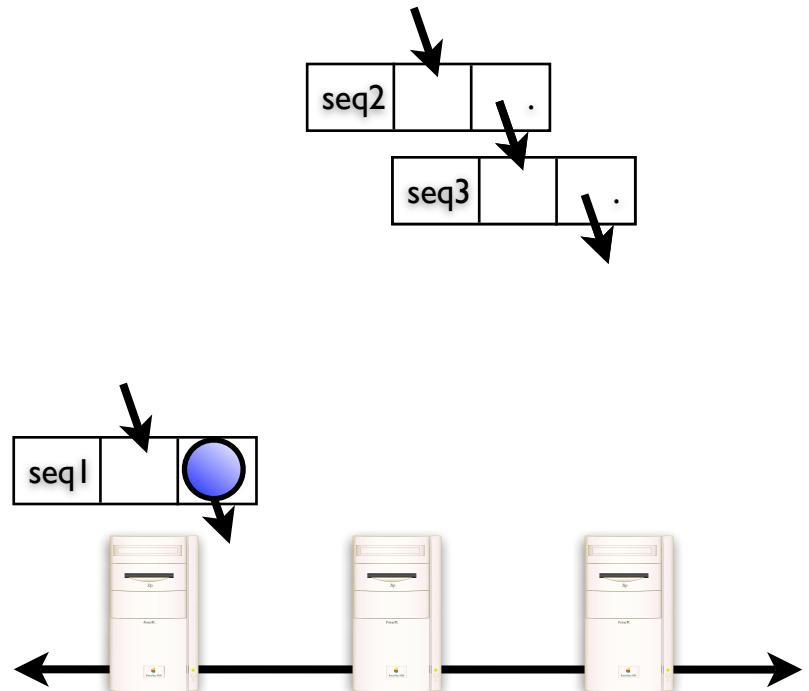
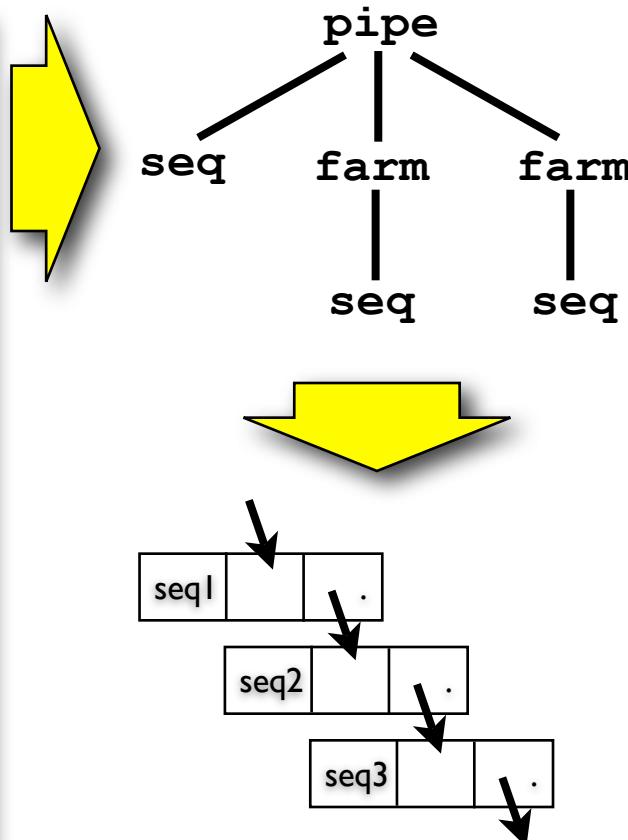
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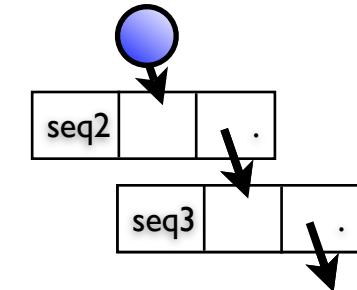
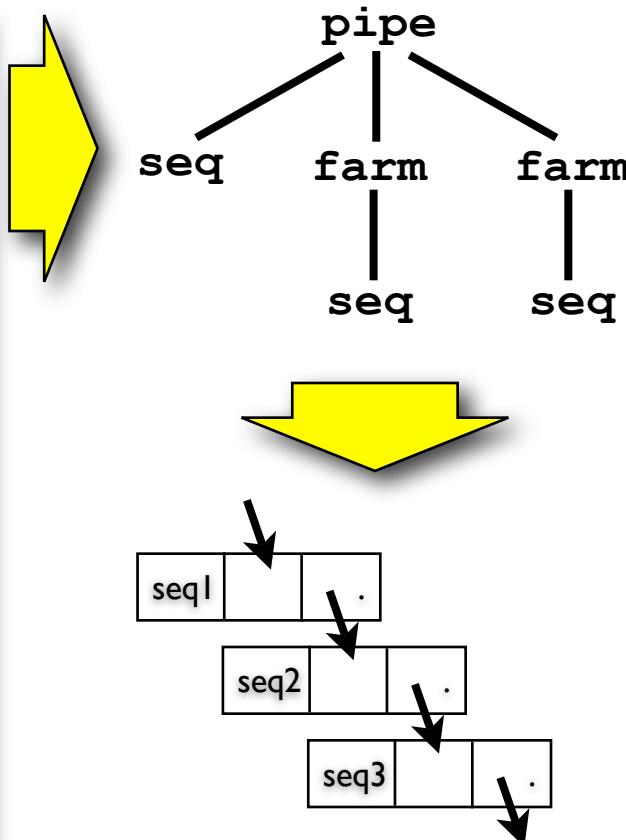
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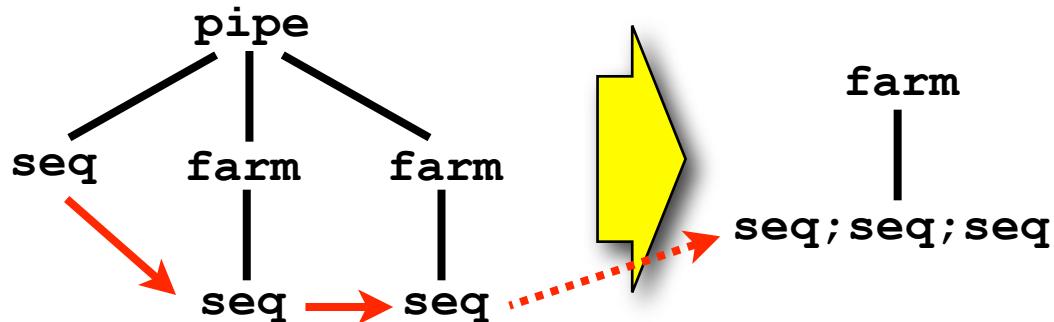
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source code



this is Lithium/muskel ('00)

# Normal form ('99)



$\text{Skel} = \text{Seq} \mid \text{farm}(\text{Skel}) \mid \text{pipe}(\text{Skel}, \text{Skel})$

- $\text{NF}(\text{Seq}) = \text{Seq}$
- $\text{NF}(\text{farm}(\text{Skel})) = \text{NF}(\text{Skel})$
- $\text{NF}(\text{pipe}(\text{Skel}', \text{Skel}'')) = \text{NF}(\text{Skel}') ; \text{NF}(\text{Skel}'')$
- $\text{normalForm}(\text{Skel}) = \text{farm}(\text{NF}(\text{Skel}))$

- + increases computational grain
  - + decreases comm overhead
  - + improves service time
  - + uses leg amount of PEs



# MDF + NF

- Locality handled by support rather than programmer
  - affinity scheduling of MDF instructions
- Grain handled by the support
  - communication clustering / multithreaded remote interpreters
- Better skeleton sets to be investigated (meta annotations vs. skeletons)



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# Generative aspects

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- Layered implementation of structured parallel programming environments
- Each layer with precise responsibilities
  - local optimizations
- Strictly hierarchical layer layout
  - avoid local optimization interference



# Layers

## Applications

Qualitative parallelism exploitation, user knowledge

## Compiler tools

Static code generation + static transforms/optimizations, system designer knowledge

## Loader / Deployment tools

Target machine dependent code choice, system designer + machine expert knowledge

## Run time system

Adaptivity, fault tolerance, QoS, ... , machine expert knowledge



# Layered systems

```
generic main()
{
    stream T1 s1;
    stream T2 s2;

    stage1(output_stream s1);
    stage2(input_stream s1,
            output_stream s2);
    stage3(input_stream s2);
}

proc s1(output_stream T1
        s) $c{ /* C code */ }c$;

stage1(output_stream T1 s)
{
    s1(output_stream s)
}

parmod stage2(input_stream
T1 sin, output_stream T2
sout) {
```

**source code**

**astCC**

**obj code:  
C++  
ASSISTlib**

**multi target  
makefiles**

**XML  
config  
file**

ready to compile  
code

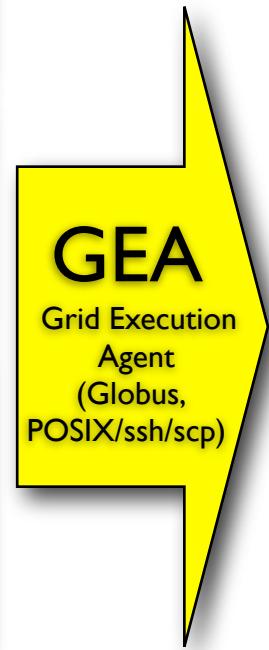
relative to diff  
target archs

sw config (libs,  
run time) + proc  
network config

**this is ASSIST ('00)**



# Layered systems



- + figures out target machine features
- + compiles proper code subsets using proper makefiles
- + deploys/stages user and system code
- + starts up run time system processes, user (derived) processes, management processes
- + stages input data
- + starts and monitors program execution
- + stages back output data

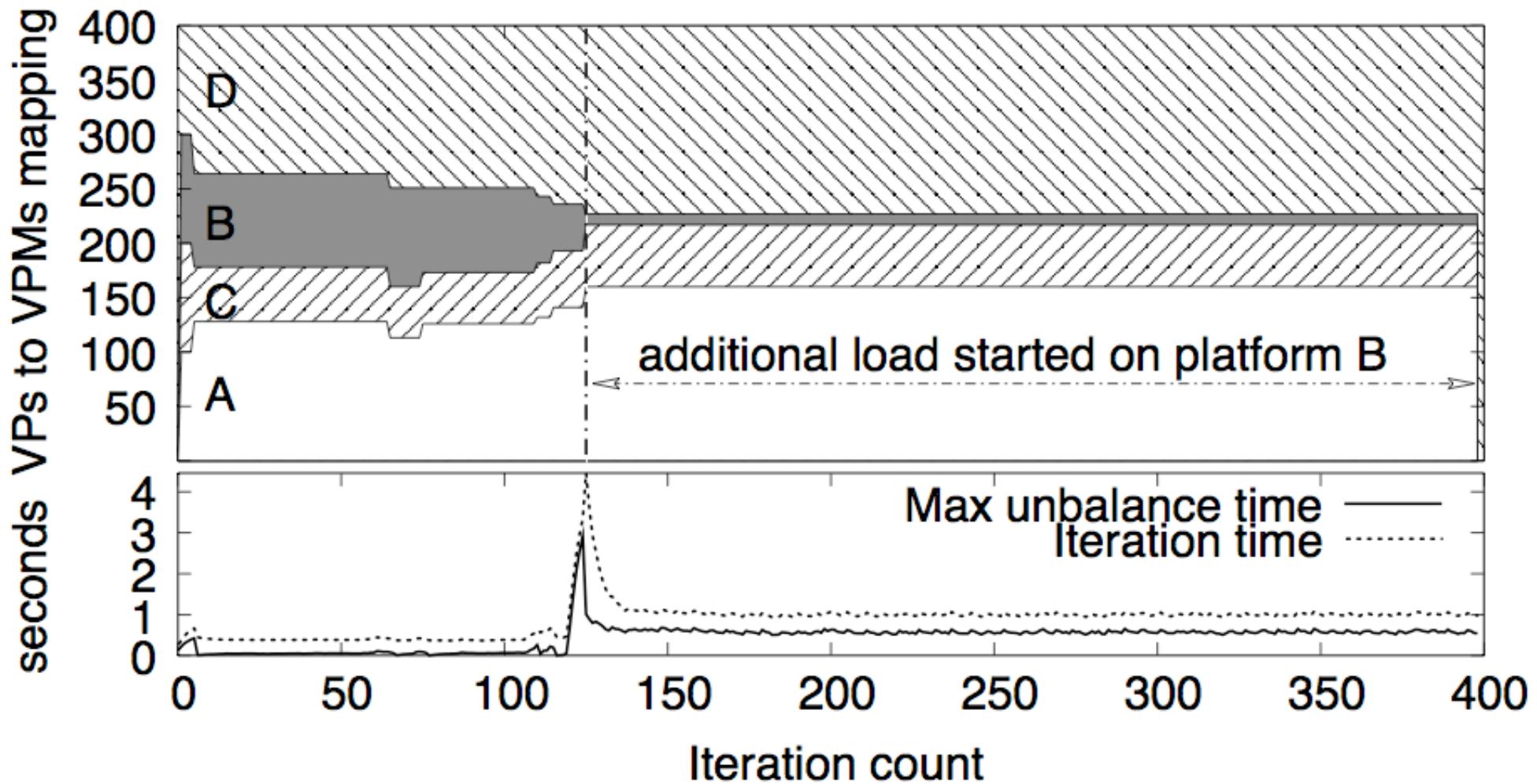


# Plus ...

- interoperability
  - compiler wraps ASSIST code to WS/CCM
- fault tolerance
  - comp. generated control code works at run time
- adaptivity
  - compiler generates manager code
  - manager adapts prog behavior at run time



# ASSIST adaptivity





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- Skeletons provide very high level information, exploitable to generate good code
- Separation of concerns is fundamental to provide efficiency
- Proper choice of compile/load/run time timing is fundamental as well



- 
- any questions ?

( [marcod@di.unipi.it](mailto:marcod@di.unipi.it) )



# References

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- P3L: Susanna Pelagatti "Structured Development of Parallel Programs"  
Taylor&Francis '98
- muskel: [www.di.unipi.it/~marcod](http://www.di.unipi.it/~marcod)
- ASSIST: [www.di.unipi.it/Assist.html](http://www.di.unipi.it/Assist.html)
- [marcod@di.unipi.it](mailto:marcod@di.unipi.it)